Smallest Explanations and Diagnoses of Rejection in Abstract Argumentation

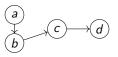
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September 16, 2020 @ KR 2020, Online

Argumentation in AI

- Active and vibrant area of modern AI research
- Central KR formalism for reasoning in abstract argumentation: argumentation frameworks (AFs) [Dung, 1995]



Explaining and Diagnosing in Abstract Argumentation

- Understanding reasons for rejection important and nontrivial
- Diagnosing why no argument is accepted [Ulbricht and Baumann, 2019]
- Explaining credulous rejection of an argument

[Saribatur et al., 2020]

What?

- **Provide complexity results** for computing smallest explanations and diagnoses of credulous rejection of a given argument
- Design declarative algorithms for practical computation
 - both argument-based and attack-based explanations and diagnoses

How?

Identify correspondences between

- minimal (smallest) explanations and (smallest) MUSes
- minimal (smallest) diagnoses and (smallest) MCSes

of propositional formulas in CNF

 $\begin{aligned} \mathsf{MUS} &= \mathsf{minimal} \text{ unsatisfiable subset} \\ \mathsf{MCS} &= \mathsf{minimal correction set} \end{aligned}$

Given an AF
$$F = (A, R)$$
, $q \in A$, $\sigma \in \{adm, stb\}$.

Definition

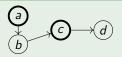
A set $A' \subseteq A$ of arguments is an **explanation** for rejecting q: q remains rejected in any sub-AF containing A'

Definition

A set $A' \subseteq A$ of arguments is a **diagnosis** of rejecting q: q becomes accepted in sub-AF where A' is removed

Example

 $\{a, c\}$ is an explanation for rejecting d



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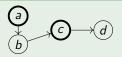
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 $\{a, c\}$ is an explanation for rejecting d



Given an AF F = (A, R), $q \in A$, $\sigma \in \{adm, stb\}$, and an integer $k \ge 0$.

Theorem

Deciding whether there exists an explanation $A' \subseteq A$ with $|A'| \leq k$ for rejecting q in F under σ is Σ_2^{p} -complete.

Consider the standard reduction from CNF to AFs. Reduce from deciding whether there is an unsatisfiable subset of size at most k. [Liberatore, 2005]

Theorem

Deciding whether there exists a diagnosis $A' \subseteq A$ with $|A'| \leq k$ of rejecting q in F under σ is **NP-complete**.

Reduce from credulous acceptance under σ .

Given an AF F = (A, R), $q \in A$, $\sigma \in \{adm, stb\}$.

 \Rightarrow Propositional formulas (with hard and soft clauses) for which

- an MUS corresponds to a minimal explanation,
- an MCS corresponds to a minimal diagnosis.

Computation of Smallest Explanations and Diagnoses

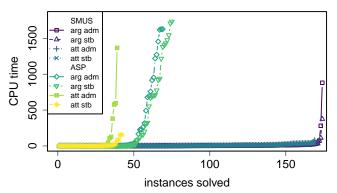
Declaratively via computing smallest MUS/MCS using

- system for extracting smallest MUS [Ignatiev et al., 2015]
- MaxSAT solver for computing smallest MCS

[Ignatiev et al., 2019]

Implementation available online in open source: https://bitbucket.org/andreasniskanen/selitae Comparison to recent ASP-based approach for computing smallest explanations

[Saribatur et al., 2020]



explanation

Paper Summary

- Complexity results for deciding small explanations and diagnoses
 - Σ_2^{p} -completeness and NP-completeness
- Algorithms for computing smallest explanations and diagnoses
 - employing smallest MUS extractors and MaxSAT solvers

Future Outlook

- Complexity of attack-based explanations and diagnoses open
- Dually: explaining and diagnosing skeptical acceptance



Dung, P. M. (1995).

On the acceptability of arguments and its fundamental role in nonmonotonic reasoning, logic programming and n-person games.

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Ignatiev, A., Morgado, A., and Marques-Silva, J. (2019). RC2: An efficient MaxSAT solver. *J. Satisf. Boolean Model. Comput.*, 11(1):53–64.



Ignatiev, A., Previti, A., Liffiton, M. H., and Marques-Silva, J. (2015). Smallest MUS extraction with minimal hitting set dualization. In *CP*, volume 9255 of *LNCS*, pages 173–182. Springer.



Liberatore, P. (2005). Redundancy in logic I: CNF propositional formulae. *Artif. Intell.*, 163(2):203–232.



Saribatur, Z. G., Wallner, J. P., and Woltran, S. (2020). Explaining non-acceptability in abstract argumentation. In *ECAI*, volume 325 of *FAIA*, pages 881–888. IOS Press.

Ulbricht, M. and Baumann, R. (2019).

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