

# Aggregating Discrete Information from Mutually Inconsistent Sources

Nir Ailon

based on joint work with Moses Charikar and Alantha Newman



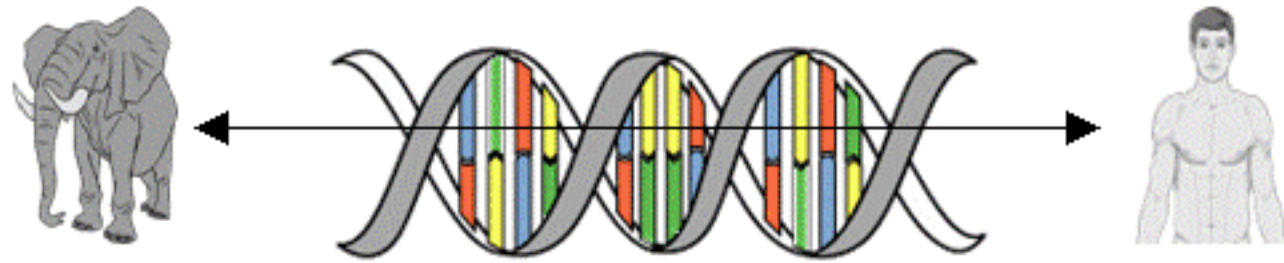
reconstructing . . . .



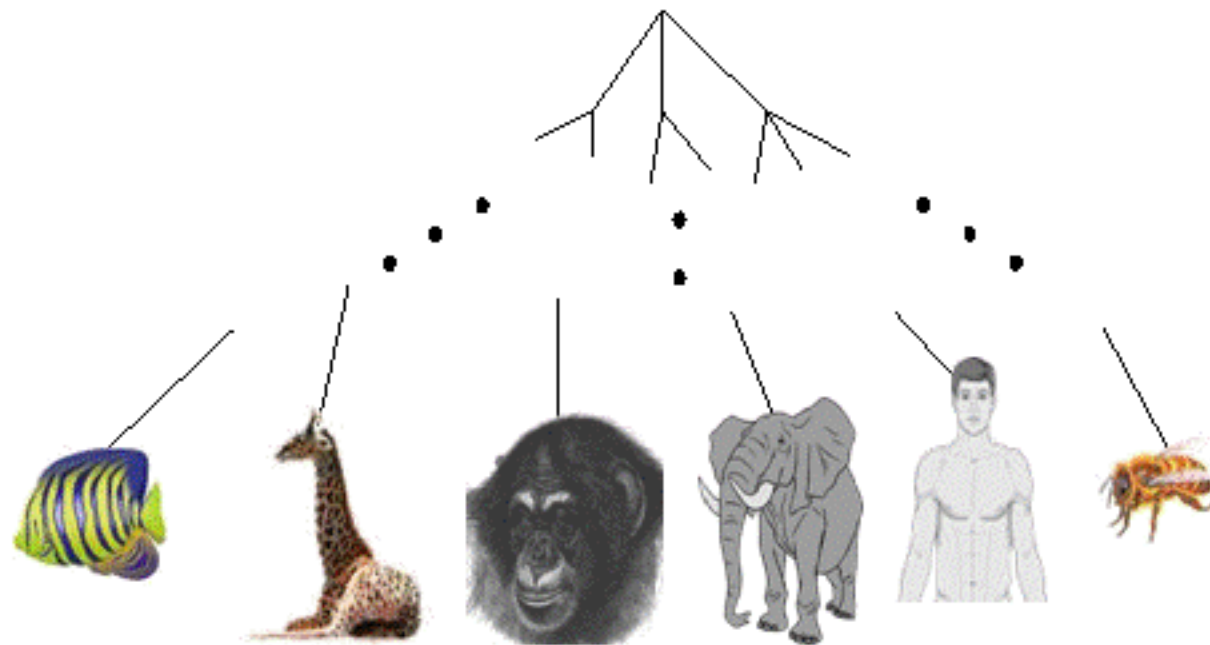


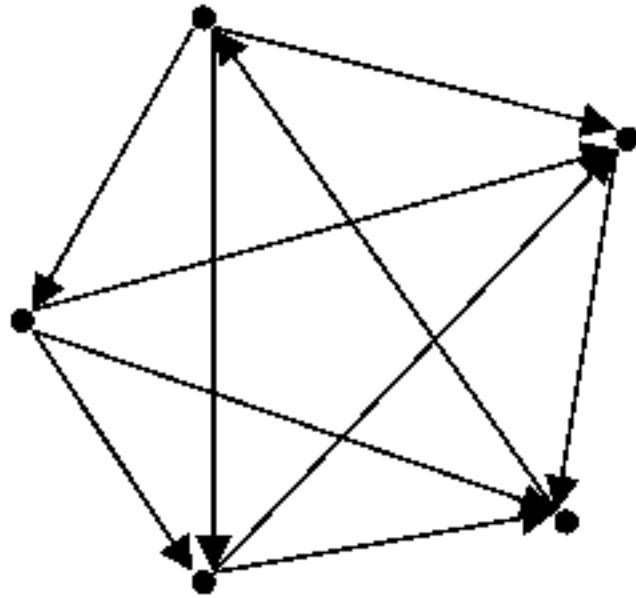
reconstructing . . . .



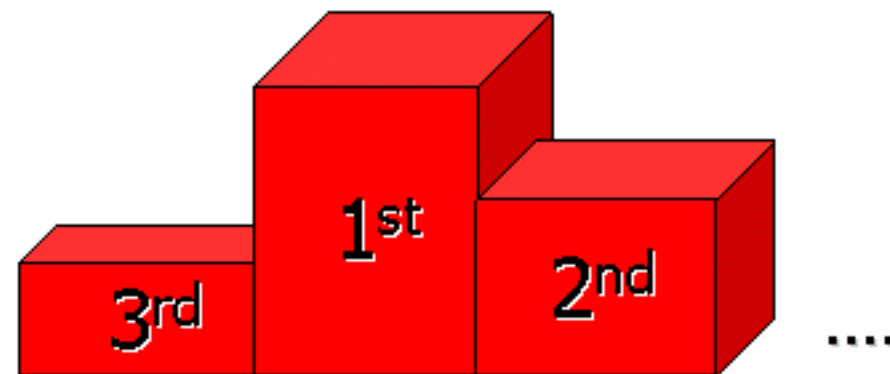


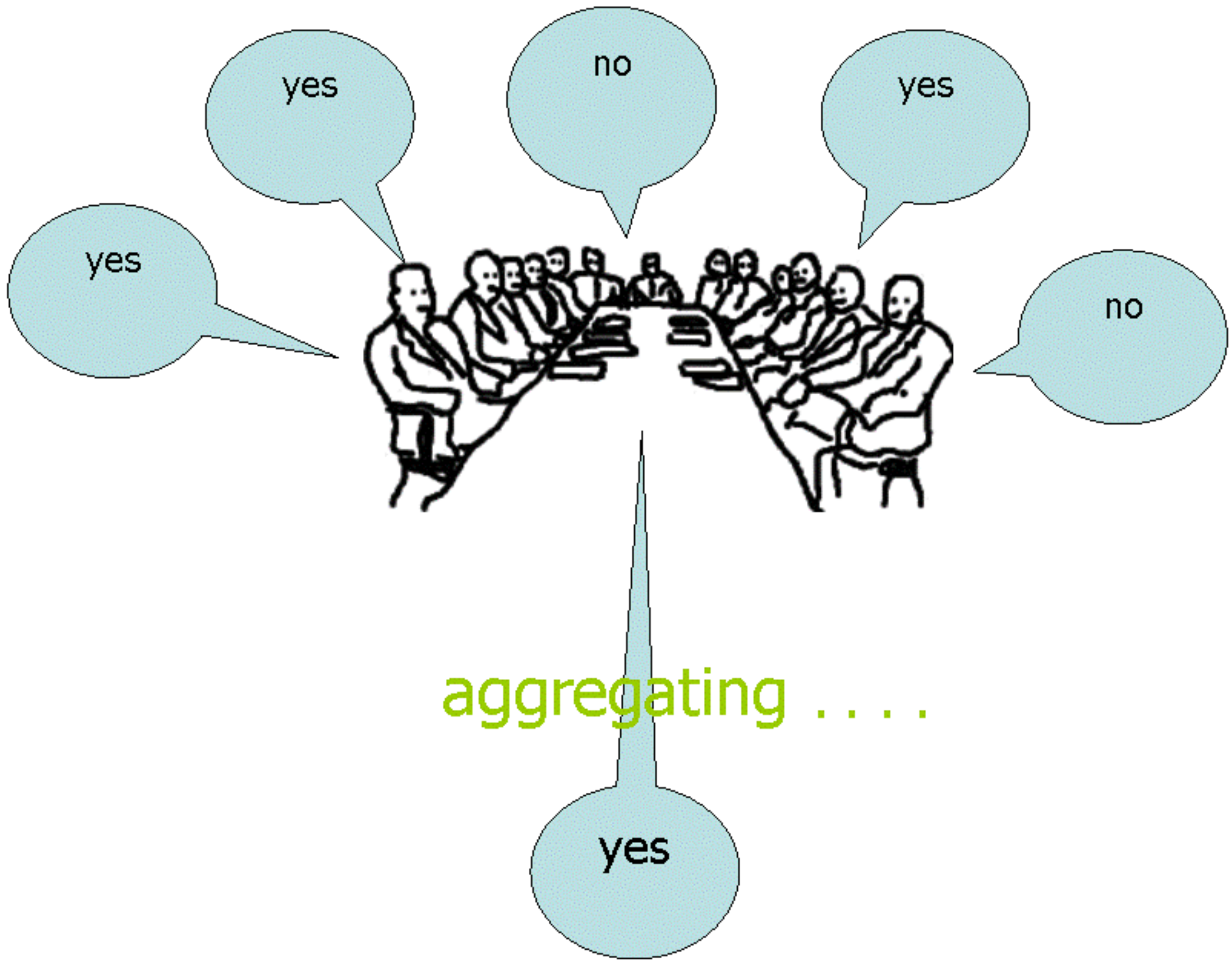
reconstructing . . . .

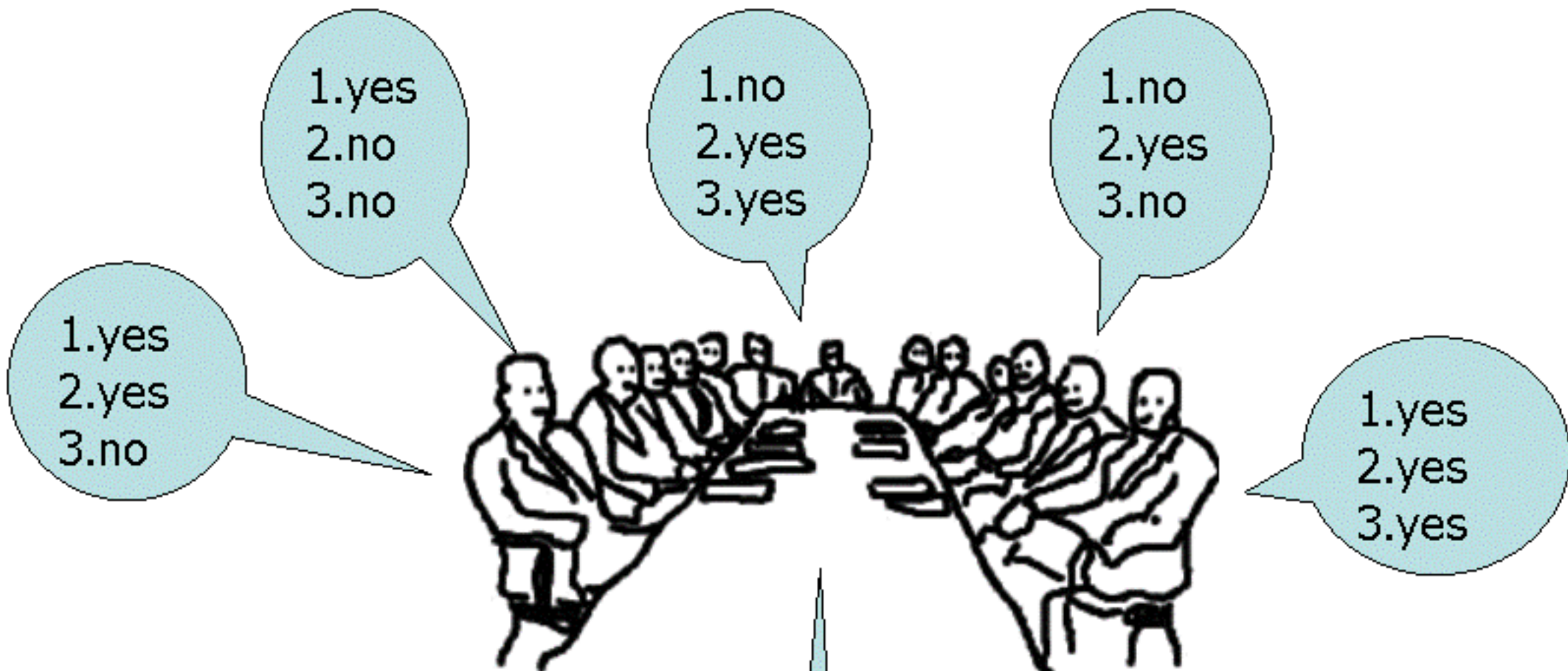




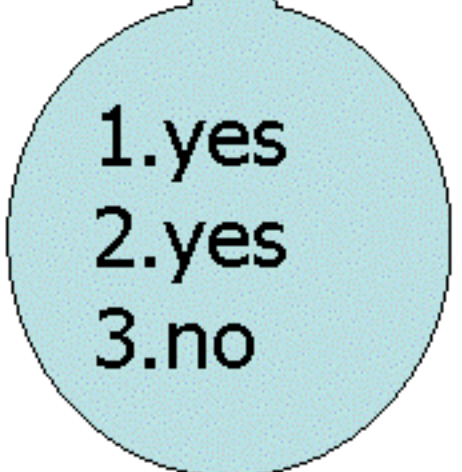
reconstructing . . . .





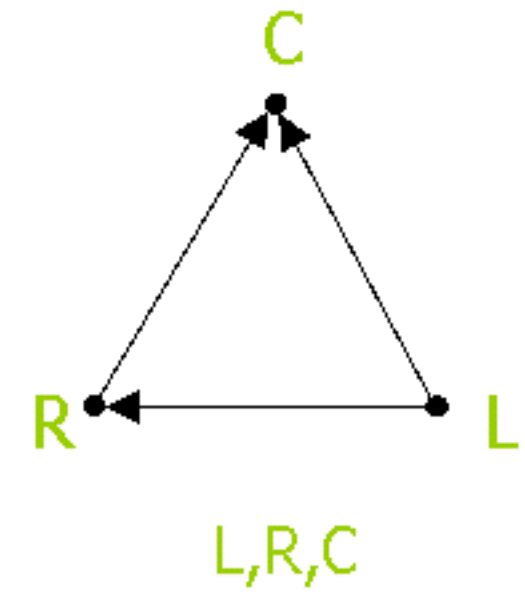
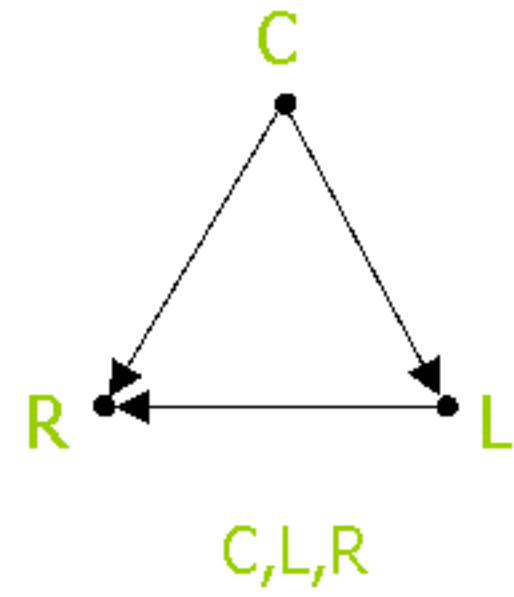
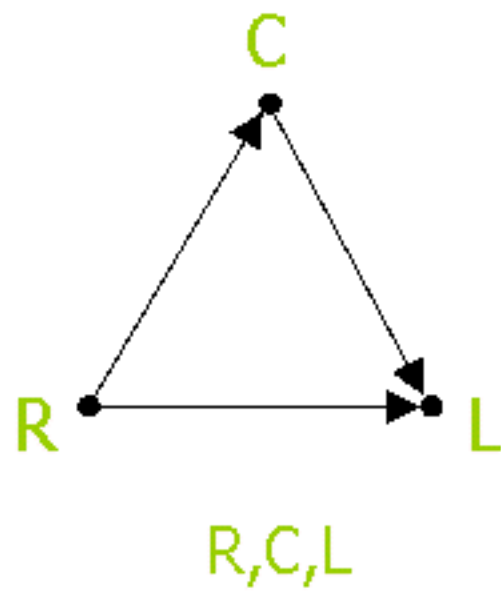


aggregating . . . .

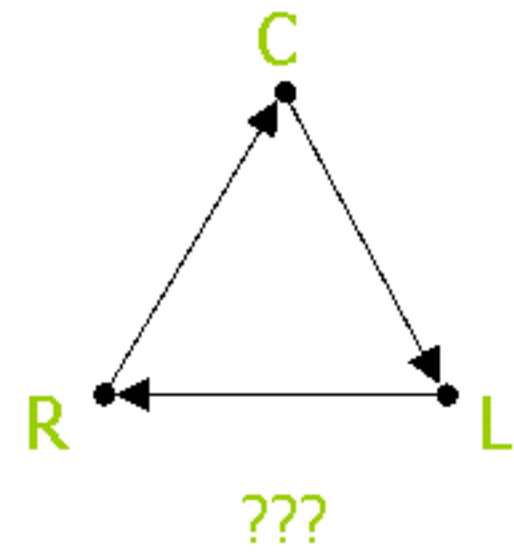


1. product **R**elease before advertising **C**ampaign?
2. advertising **C**ampaign before massive **L**ayoffs?
3. massive **L**ayoffs before product **R**elease?





aggregating . . . .



# Ranking Applications

- Information retrieval
- Voting (sports, politics...)
- Paleontology
- Finance
- Movie/hotel/restaurant rating
- Experimental psychology
- Clustering
- Collaborative filtering

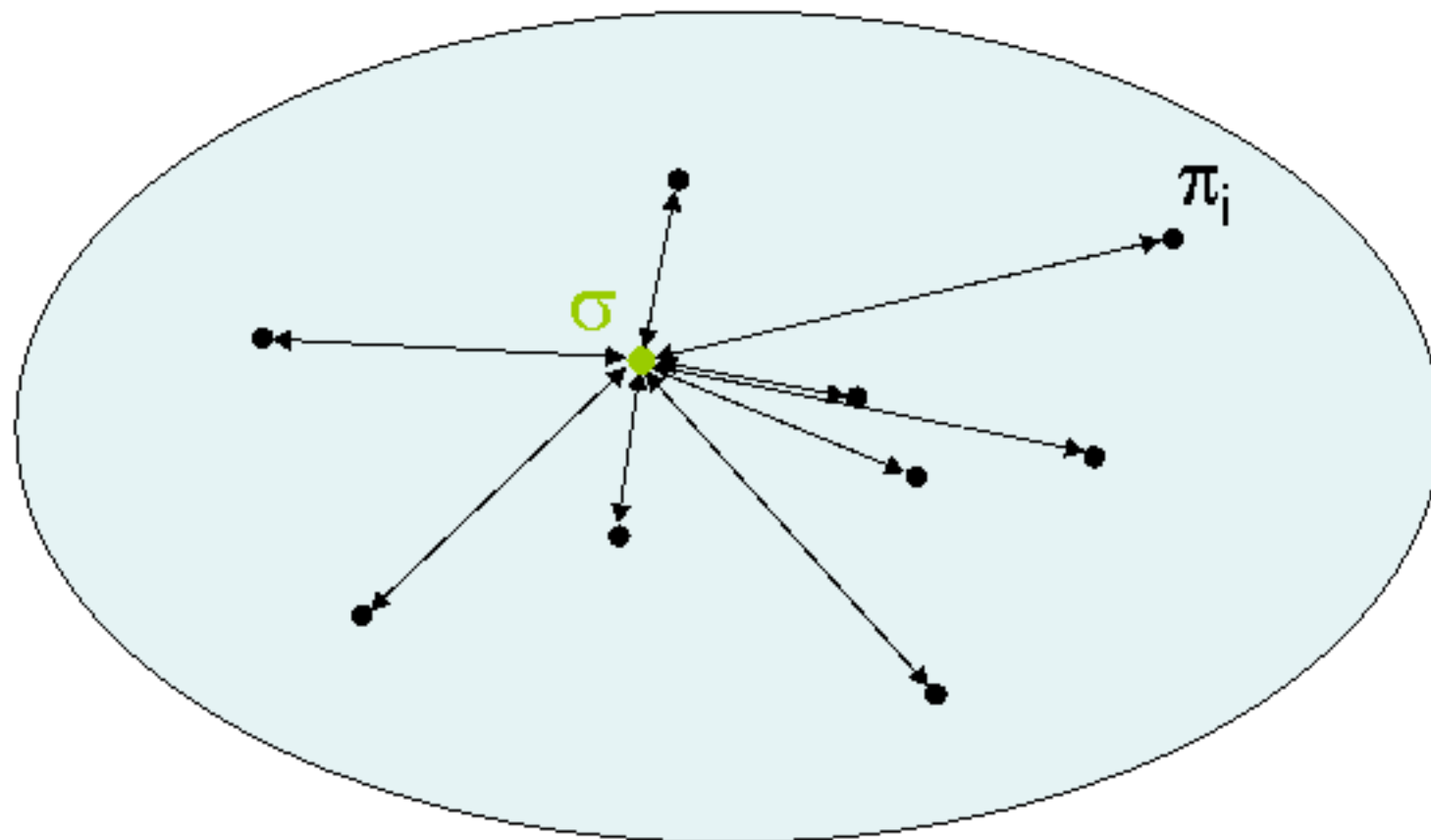
## Rank Aggregation: *(Kemeny-Young, minimization)*

$$\text{OPT} = \min_{\sigma} \text{cost}(\sigma) = \min_{\sigma} \sum_{i=1..k} \langle \sigma, \pi_i \rangle / k$$

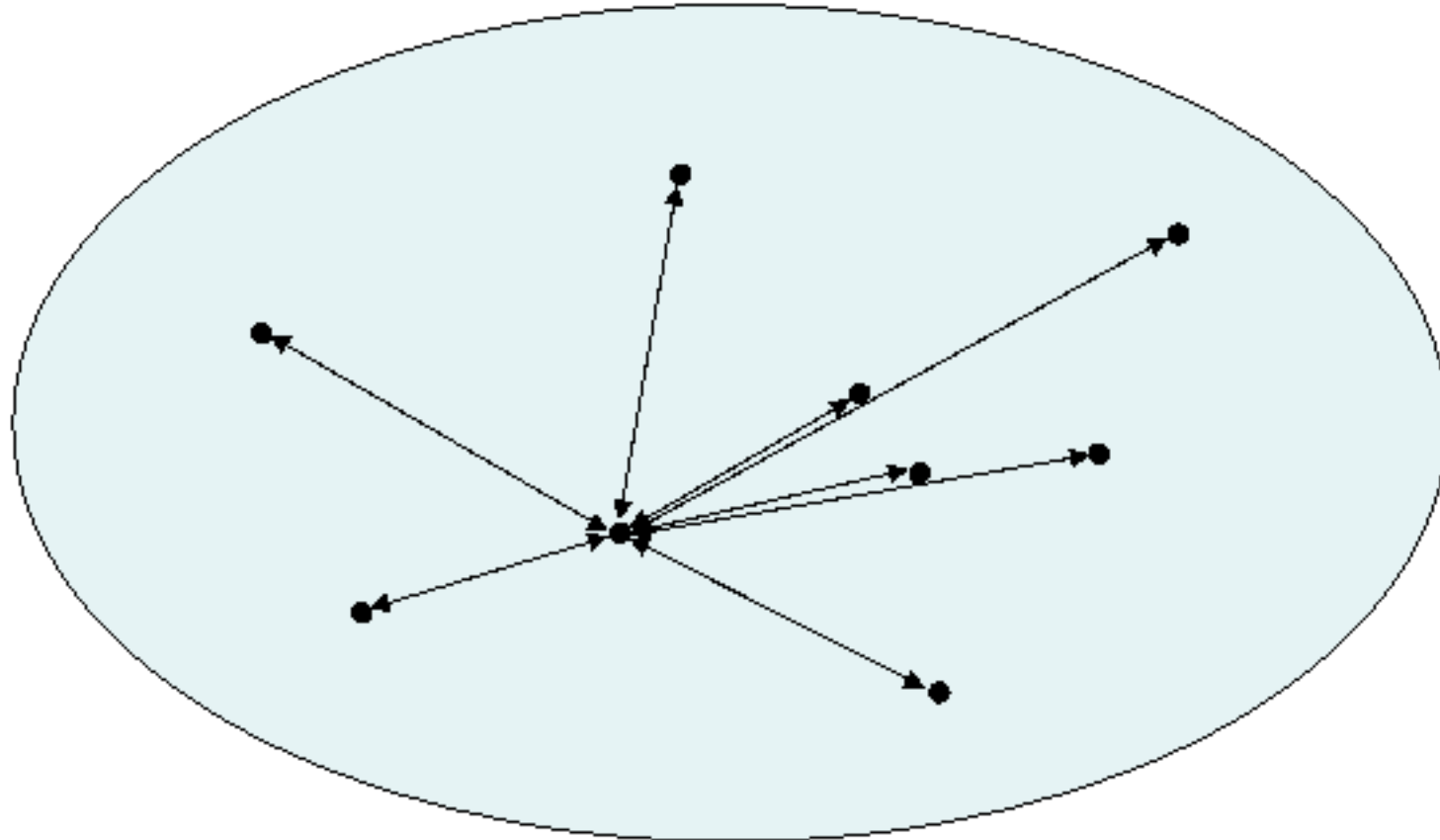
$\langle \sigma, \pi \rangle$  = inversion distance

e.g.  $\langle \text{ABC}, \text{CAB} \rangle = 2$

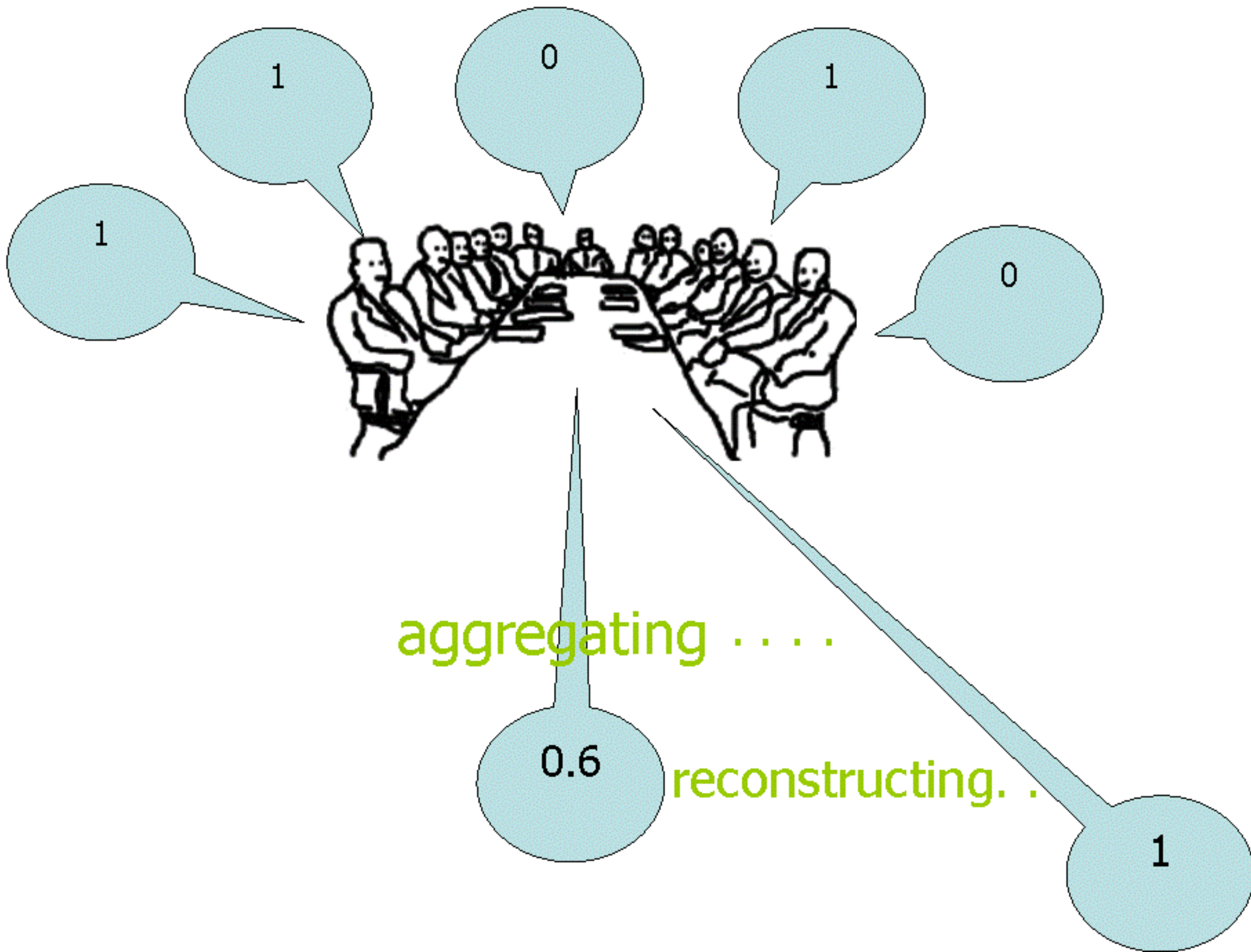
$\langle \cdot, \cdot \rangle$  metric on permutations



$$\min_j \text{cost}(\pi_j) \leq 2 \text{OPT}$$
$$\mathbb{E}_{j \in [k]} [\text{cost}(\pi_j)] \leq 2 \text{OPT}$$



Algorithm RandomChoice: 2-approx



# How to Take Average of Permutations?

write  $\pi$  as vector

score based:  $n$  coordinates

$$\pi(u) = \text{rank of } u \in V$$

pairwise based:  $n^2$  coordinates

$$\pi(u, v) = 1 \text{ if } u <_{\pi} v, 0 \text{ otherwise}$$

$$\langle \pi, \sigma \rangle = \sum \pi(u, v) \sigma(v, u)$$

$$\hat{\pi} = \sum_{i=1..k} \pi_i / k$$

$$\hat{\pi}(u, v) = \#\{i: u <_{\pi_i} v\} / k$$

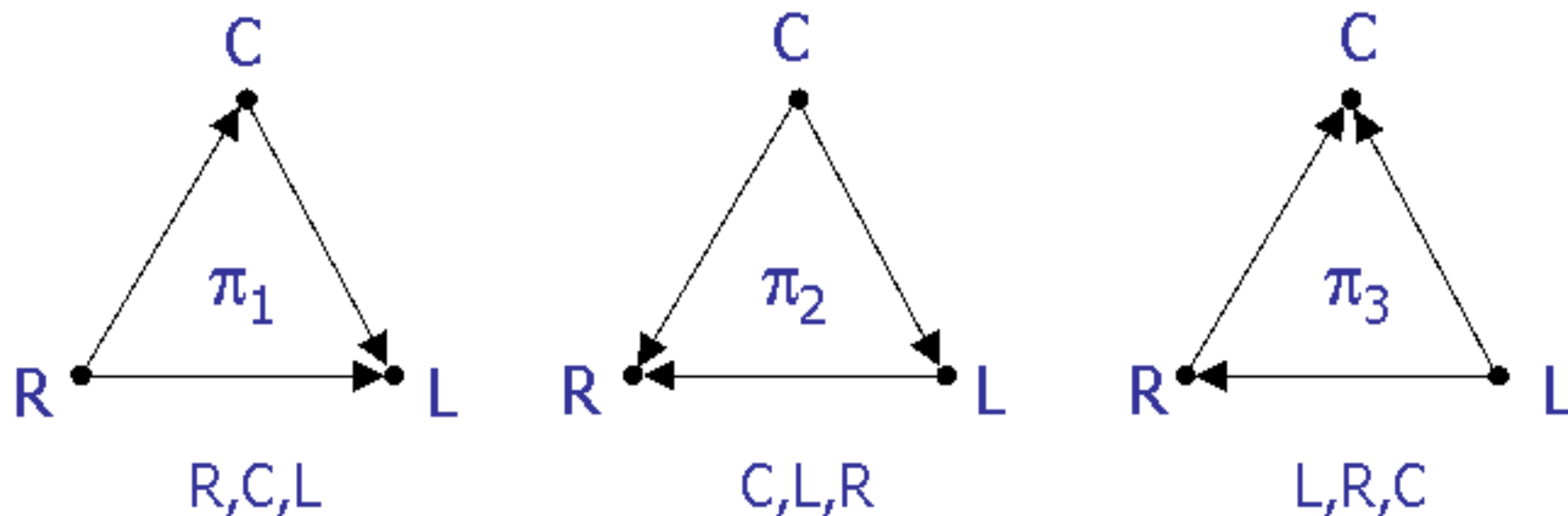
$$\text{cost}(\sigma) = \langle \sigma, \hat{\pi} \rangle$$

$$= \sum_{u <_{\sigma} v} \hat{\pi}(v, u)$$

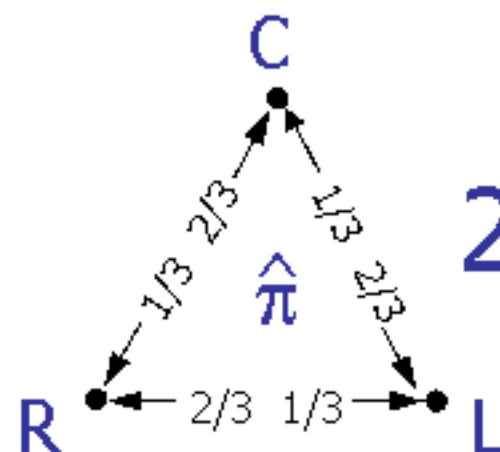
$$\hat{\pi} = \sum_{i=1..k} \pi_i / k$$

$$\hat{\pi}(u,v) = \#\{i: u <_{\pi_i} v\} / k$$

$$\text{cost}(\sigma) = \langle \sigma, \hat{\pi} \rangle$$



1. aggregating . . .



2. reconstructing . . . .

## Minimum Feedback Arc-Set (min-FAS):

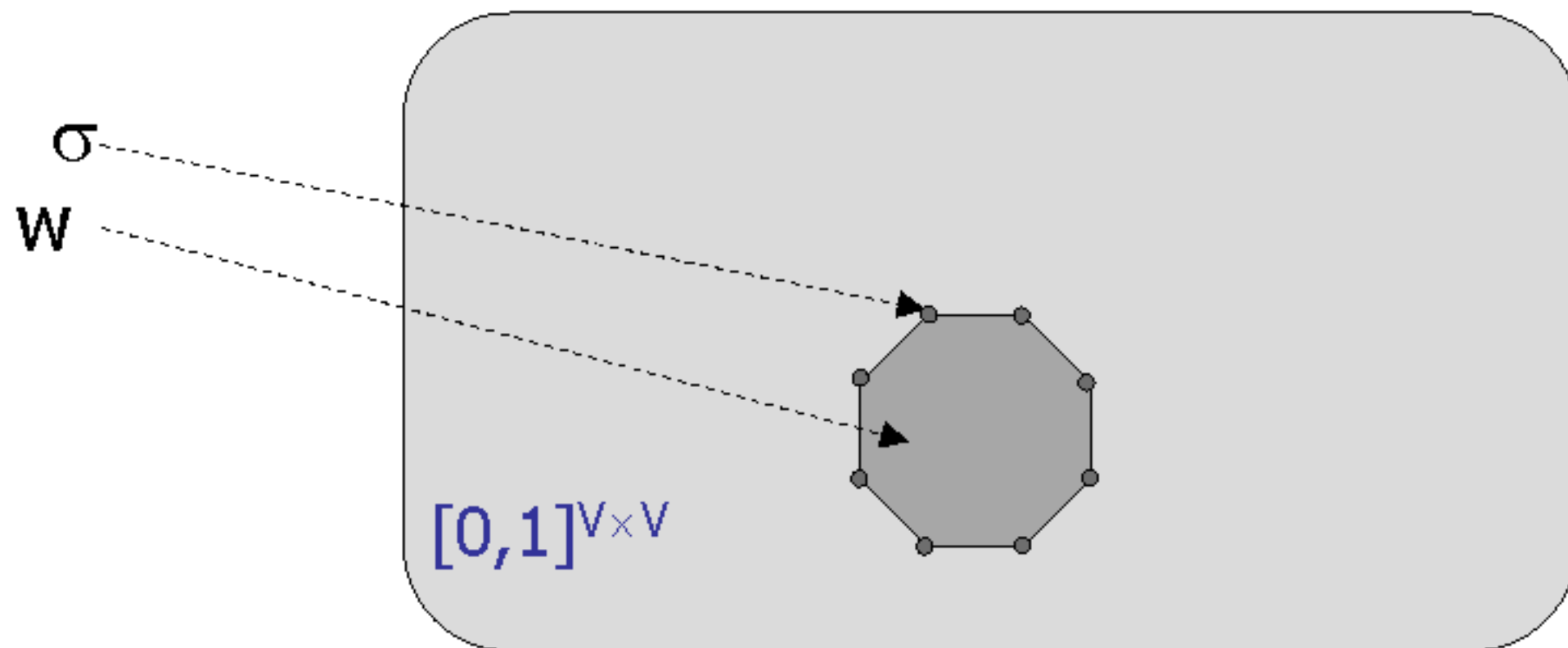
given:

set  $V$  of ( $n$  vertices)

$w: V \times V \rightarrow \mathbb{R}^+$

minimize  $\text{cost}(\sigma) = \langle \sigma, w \rangle$  over permutations  $\sigma$

Rank Aggregation:  $w = \hat{\pi} \in \text{conv}(\text{permutations})$





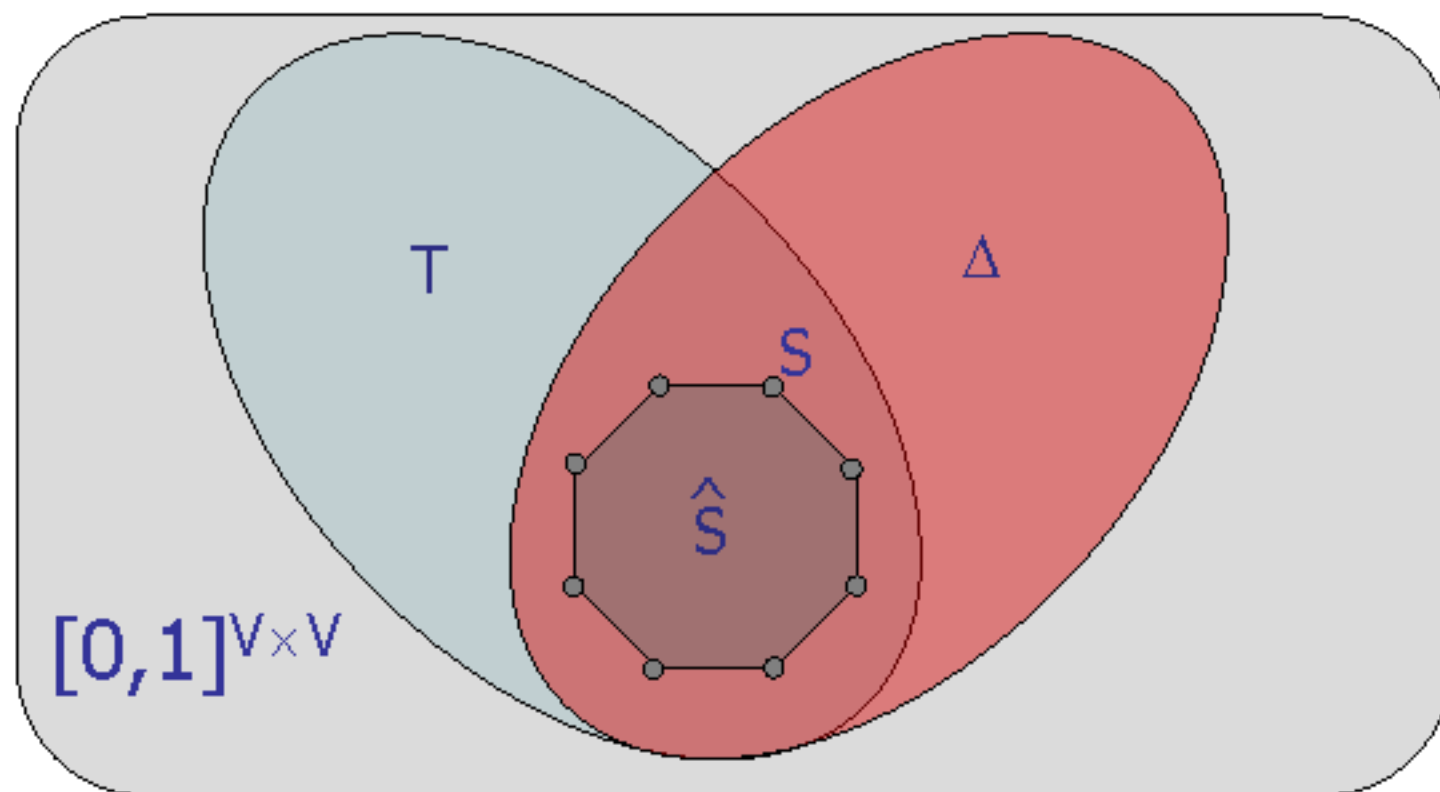
$S$ : permutations on  $V$

$\hat{S}$ : convex closure of  $S$

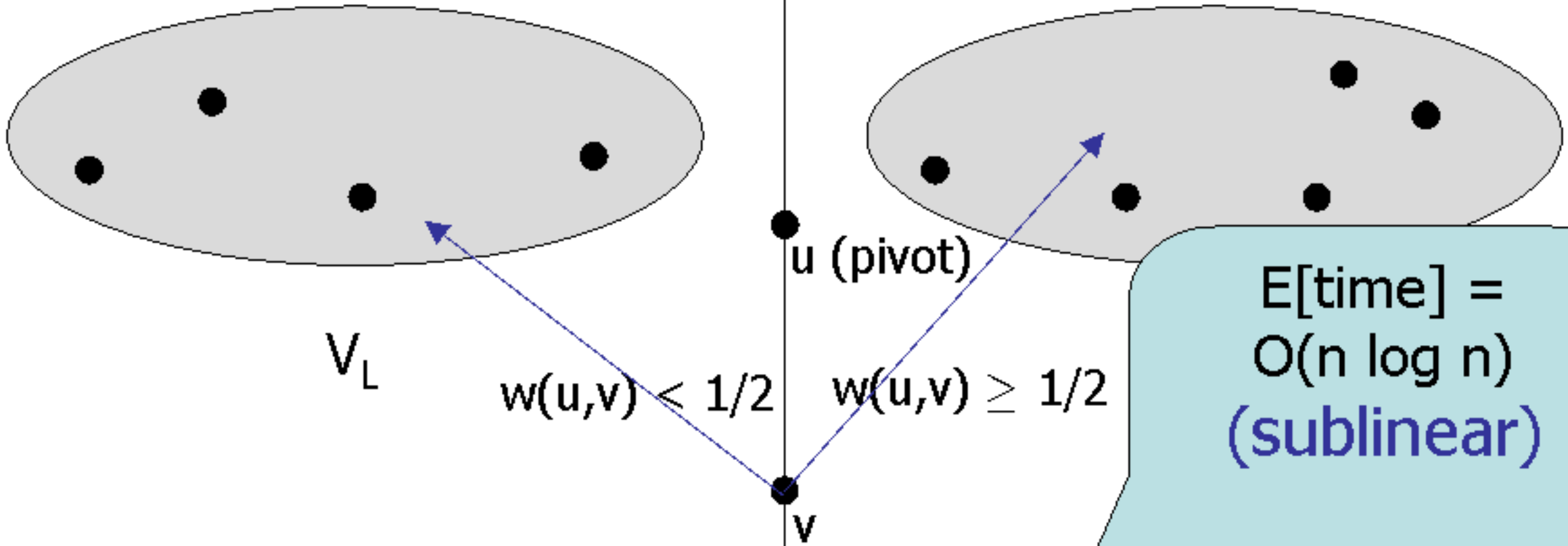
$T$ :  $\{w: w(u,v) + w(v,u) = 1 \ \forall \ u,v\}$

$\Delta$ :  $\{w: w(u,v) \leq w(u,y) + w(y,v) \ \forall \ u,v,y\}$

$S = T \cap \Delta \cap \{0,1\}^{V \times V}$



$$\text{KwikSort}(V) = \text{KwikSort}(V_L), u, \text{KwikSort}(V_R)$$



$E[\text{time}] = O(n \log n)$   
**(sublinear)**  
[GMPU 06]

(combinatorial) KwikSort

## (Combinatorial) KwikSort

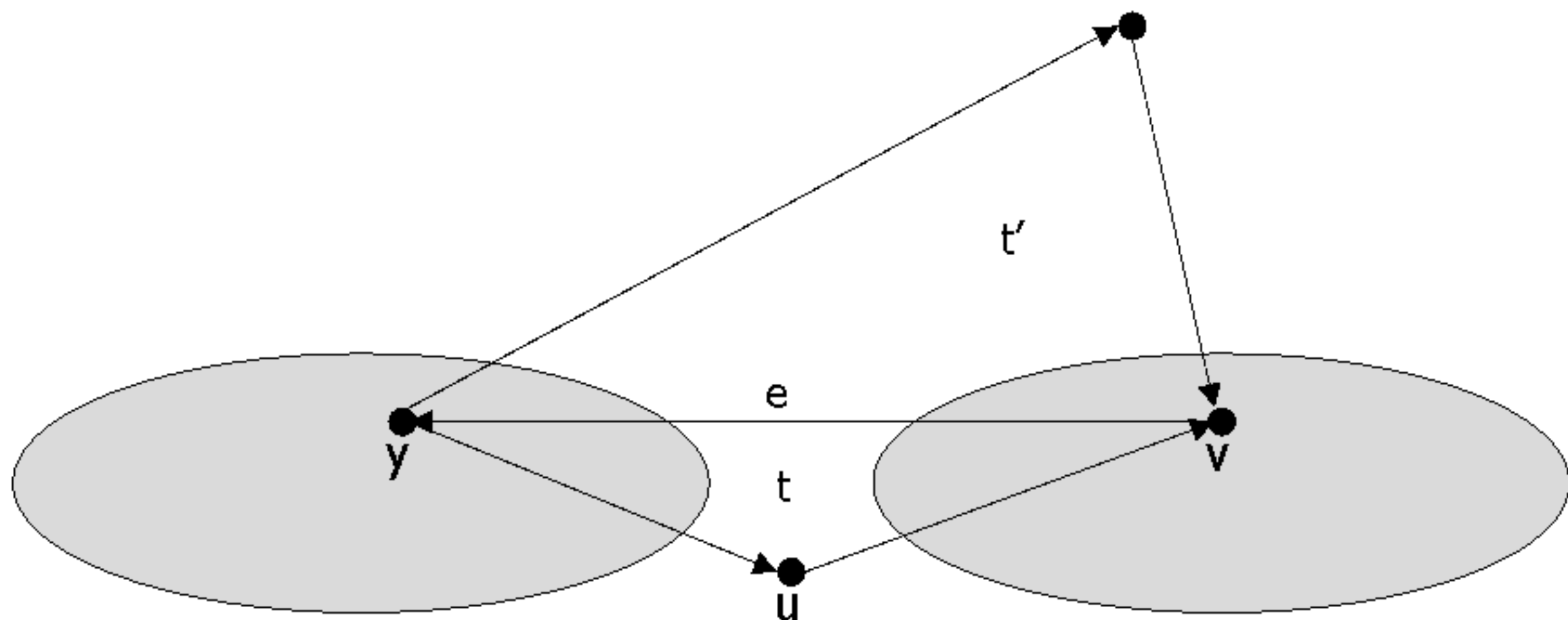
input $w$	approx	previous	hardness
$[0,1]^{V \times V}$	?	$\log n \log \log n$	Max-SNP-Hard
$T$	5	$\log n \log \log n$	NP-Hard
$T \cap \{0,1\}^{V \times V}$	3 (t'nament)	$\log n \log \log n$	NP-Hard
$\Delta$	2	$\log n \log \log n$	NP-Hard
Rank Agg	2	2	NP-Hard

best of two: 11/7

$T \cap \{0,1\}^{V \times V}$

3 (t'nament)

Proof:



$A_t$ : "triangle t charged"     $B_e$ : "edge e flipped"

$$E[\text{cost}] = \sum_t \Pr[A_t]$$

Edge e can be charged to only one triangle t

$$\Rightarrow (A_t \cap B_e) \cap (A_{t'} \cap B_e) = \emptyset, \text{ but } \Pr[A_t \cap B_e] = \Pr[A_t]/3$$

$$\Rightarrow \forall e \sum_{t: e \in t} \Pr[A_t]/3 \leq 1$$

$$\Rightarrow \sum_t \Pr[A_t]/3 \leq \text{OPT (LP duality – see next slide)}$$

$A_t$ : "triangle  $t$  charged"     $B_e$ ="edge  $e$  flipped"

$$E[\text{cost}] = \sum t \Pr[A_t]$$

Edge  $e$  can be charged to only one triangle  $t$

$$\Rightarrow (A_t \cap B_e) \cap (A_{t'} \cap B_e) = \emptyset, \text{ but } \Pr[A_t \cap B_e] = \Pr[A_t]/3$$

$$\Rightarrow \forall e \sum_{t: e \in t} \Pr[A_t]/3 \leq 1$$

$$\Rightarrow \sum t \Pr[A_t]/3 \leq \text{OPT (LP duality)}$$

Primal: (hitting LP)

minimize total value of weights assigned to edges  
such that weights incident to each directed triangle  
sum up to at least 1

Dual: (packing LP)

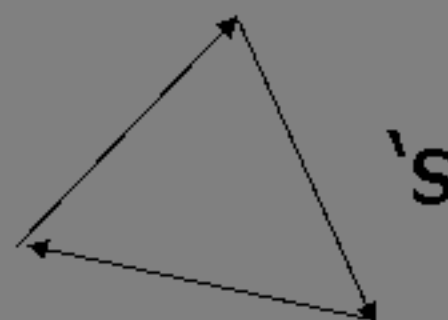
maximize total value of weights assigned to directed  
triangles such that weights incident to each edge  
sum up to at most 1

# How to Prove Other Approx Factors?

(Combinatorial) KwikSort

input w	approx	previous	hardness
$[0,1]^{V \times V}$	?	$\log n \log \log n$	Max-SNP-Hard
T			NP-Hard
$T \cap \{0,1\}^{V \times V}$			NP-Hard
$\Delta$			NP-Hard
Rank Agg			NP-Hard

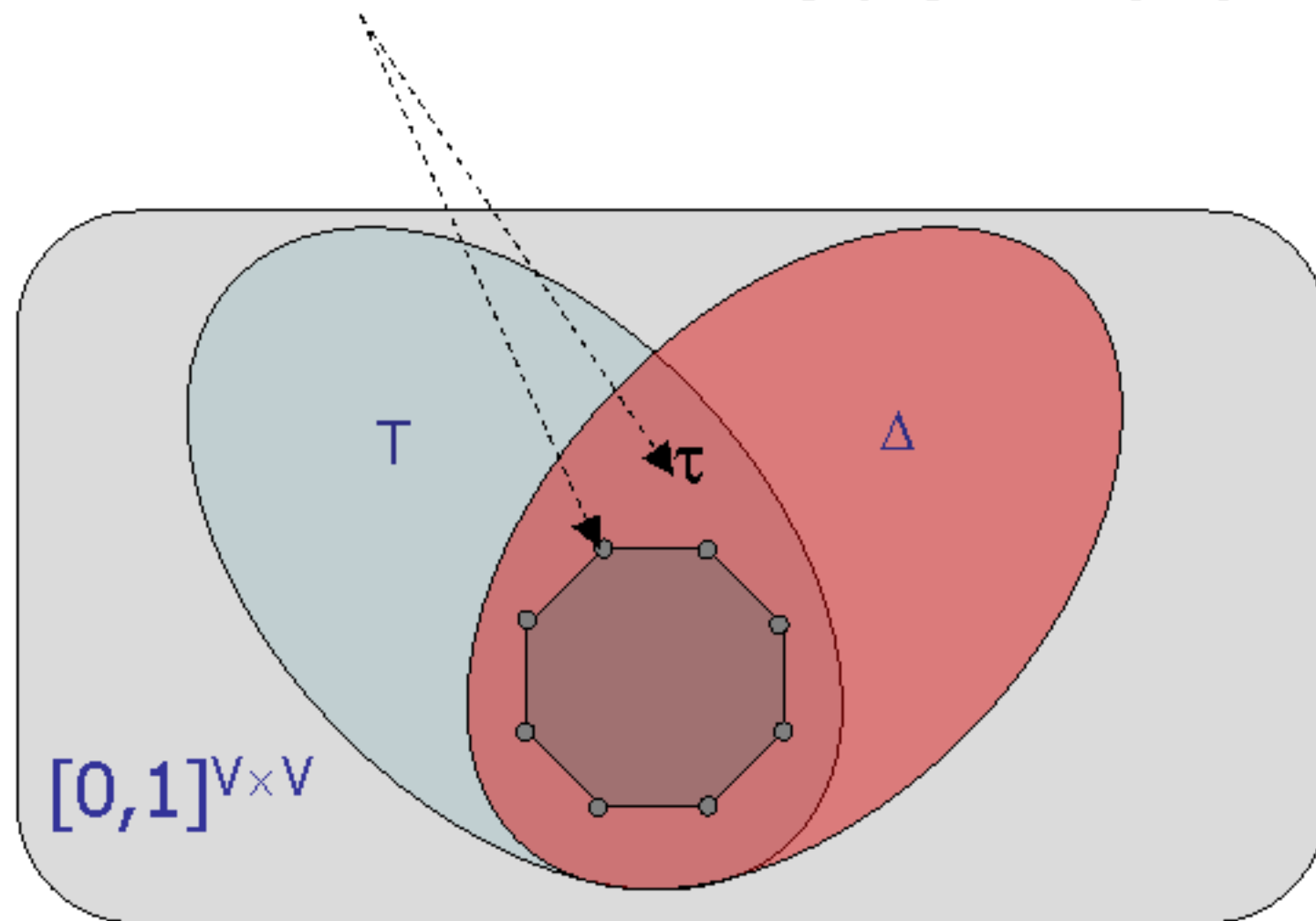
Look at



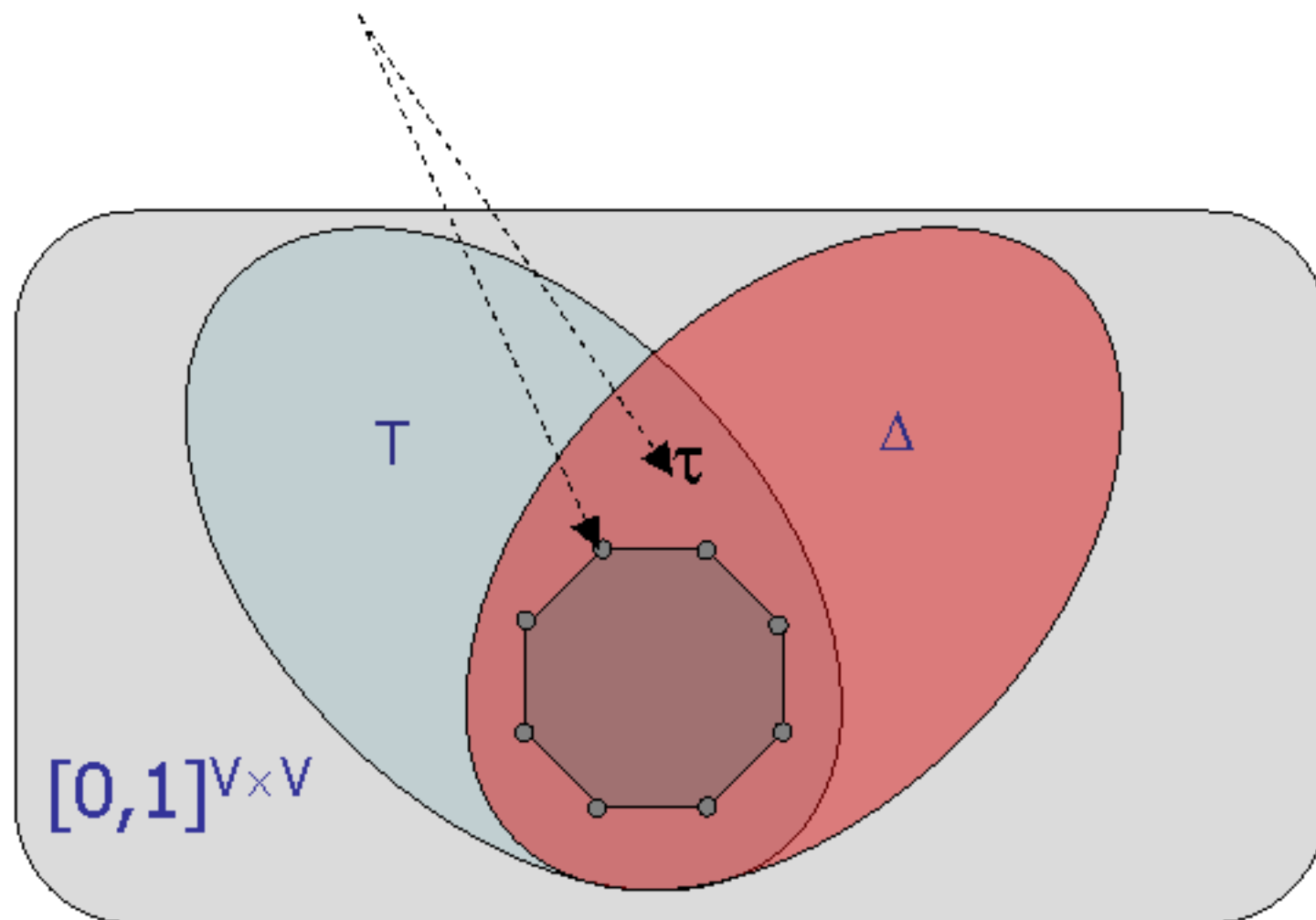
best of two: 11/7

Any structural/statistical information on weights around triangles in the input could translate to improved approximation

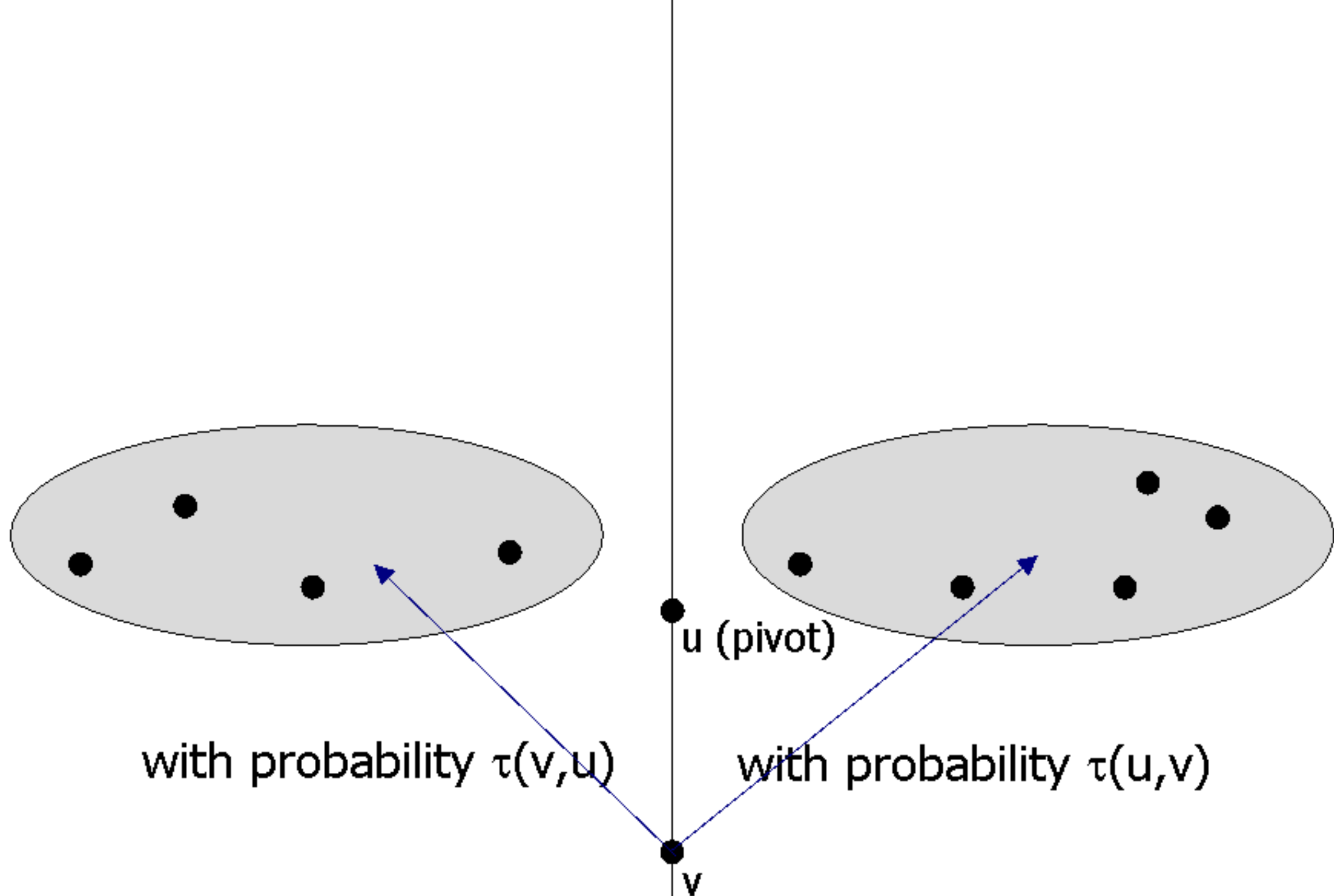
optimize over  $T \cap \Delta \cap \{0,1\}^{V \times V}$  (IP)



optimize over  $\mathcal{T} \cap \Delta$  (LP)







**(LP-based) KwikSort**

## (LP-based) KwikSort

input $w$	approx	previous	hardness
$[0,1]^{V \times V}$	?	$\log n \log \log n$	Max-SNP-Hard
$T$	2.5	5	NP-Hard
$T \cap \{0,1\}^{V \times V}$	2.5 (t'nament)	3	NP-Hard
$\Delta \cap T$	2	2	NP-Hard
Rank Agg	2	2	NP-Hard

best of two:  $4/3$     $11/7$

Ranking with Ties

a.k.a.

Bucket Order

a.k.a.

Partial Ranking



Confusing!

[Controversy](#) page. For questions or comments, please contact [Nancy O'Brien](#).

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## BUSINESS SCHOOLS

### [Asia's Best MBA Schools 2000](#)

[<http://www.asiaweek.com/asiaweek/features/mba/index.html>]

**Methodology:** This site does not offers an explanation of its rankings.

(Please note: we have recently been informed that Asiaweek will no longer be publishing their annual rankings. We will continue to link to the site as long as it is relevant.)

Asiaweek offers its assessment of the top MBA programs in the Asia-Pacific region. The ratings are based on six categories: academic resources, student selectivity, faculty resources, university connections, grade output, and other resources such as library spending and technology. Separate tables are provided for the best full-time, part-time, executive, and distance programs as well as for graduate salaries.

### [B-School Net](#)

[<http://www.b-school-net.de/>]

**Methodology:** This site offers an explanation of its rankings on the opening page.

This site, in German with English in places, offers recent rankings of Business Schools and leading MBA programs in Germany. There are several rankings including one based on a survey of 13,000 students focusing on several categories including: Quality of professors in teaching; Support for students by professors; Quality and actuality of library; Quality and availability of IT; Amount and quality of student activities; Usefulness in practice; Internationality; and Cooperation of school with corporations. There are also links to further information.

### [Business School Research Rankings](#)

[<http://www.bus.indiana.edu/ardennis/rankings/>]

**Methodology:** This site offers an explanation of its rankings on the page titled [Study Background](#).

Business School Research Rankings is intended to supplement a study of U.S. business school research performance published in the December issue of the *Academy of Management Journal*. The schools are

$V_{300}, V_{250}, V_{1845}, \dots, V_4, V_{50}, V_{1300}$

top-k  
aggregation  
( $k=3$ )  
NP-Hard  
for  $k \geq 2$



ties due to infeasibility

ranking of 2000 employees

$V = \{v_1, \dots, v_{2000}\}$



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Company Name	Symbol	Analyst	Rating
Allergan	AGN	Lazard Capital	Buy
AMIS Holdings	AMIS	CIBC Wld Mkts	Sector Outperform
Javelin Pharmaceutic	JAV	JP Morgan	Overweight
MasterCard	MA	Bear Stearns	Outperform
MetroPCS	PCS	UBS	Buy
MetroPCS	PCS	Bear Stearns	Outperform
MetroPCS	PCS	Banc of America Sec	Buy
MetroPCS	PCS	Wachovia	Outperform
Panacos Pharma	PANC	Punk, Ziegel & Co	Buy
Sirona Dental Systems	SIRO	Banc of America Sec	Neutral

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Go Bookmarks 44 blocked Check Look for Map AutoFill Send to hotel reviews

Review: [New York - New York Hotel and Casino](#), Las Vegas, Nevada

4.0 [New York, New York. Timeless!!!!](#)

April 03, 2006 *A TripAdvisor Member*,

Minneapolis, MN

... to Vegas and loved New York, New York. Our room ... will always stay at New York, New York. Many of the other casinos seemed dated or were way too expensive, but New York is timeless ...

Review: [New York City](#), New York

5.0 [I love New York](#)

January 05, 2006 *A TripAdvisor Member*,

Columbus, OH

Good place to visit, bad place to live.

Review: [New York - New York Hotel and Casino](#), Las Vegas, Nevada

1.0 [New York New York Hotel](#)

February 01, 2005 *A TripAdvisor Member*, Los

Angeles

We stayed in New York New York hotel bcos ... is looking for staying in New York New York bcos they are offering 2 SPA (worth of 40\$?????) you better change your mind

Review: [Ramada Plaza & Inn - New Yorker Hotel](#), New York City, New York

5.0 [New York](#)

March 26, 2003 *A TripAdvisor Member*

... is Madisson Square Garden and Times Square was 15 minutes walk away. I WILL be returning to new york and when I do I will be staying in the New Yorker again

[See the next 20 results](#)

minimize  
 $\langle \sigma, \hat{\pi} \rangle$   
 over  $\sigma \in S$

p-rating aggregation  
 (p=3)  
 in P for  $p \leq 2$



$\pi_i$

$v_1$	😊	😊	
$v_2$	😊	😊	😊
$v_3$	😊		
$v_4$	😊	😊	
$v_5$	😊	😊	
$v_6$	😊	😊	😊
$v_7$	😊		
$v_8$	😊		
$v_9$	😊	😊	😊
$v_{10}$	😊	😊	😊

$\pi_i \in T$   
 $\pi_i \in \Delta$

ties due to coarseness

ranking of 10 executives

$$V = \{v_1, \dots, v_{10}\}$$

$$\begin{aligned} \pi_i(v_{10}, v_1) &= 1 \\ \pi_i(v_1, v_{10}) &= 0 \\ \pi_i(v_5, v_4) &= \pi_i(v_4, v_5) = 0 \end{aligned}$$

$$\begin{aligned} v_{10} &< v_1 \\ v_5 &\equiv v_4 \end{aligned}$$



minimize  
 $\langle \sigma, \hat{\pi} \rangle$   
over  $\sigma \in S$

e.g.  $\langle ABC, [CA]B \rangle = 1$

If some voter ties  $u, v$  together  
no price incurred for  $u, v$

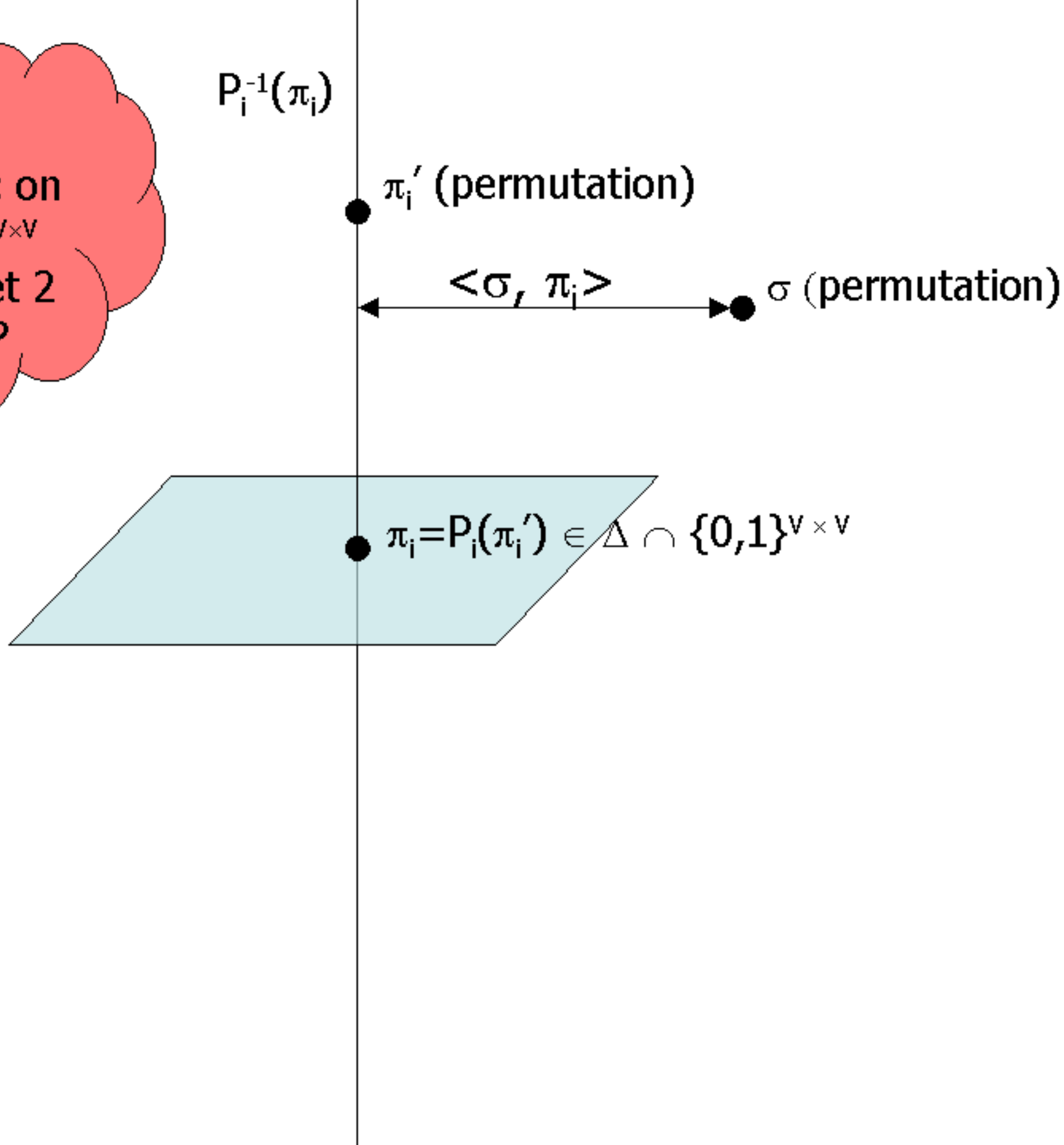
## Why not pay for ties?

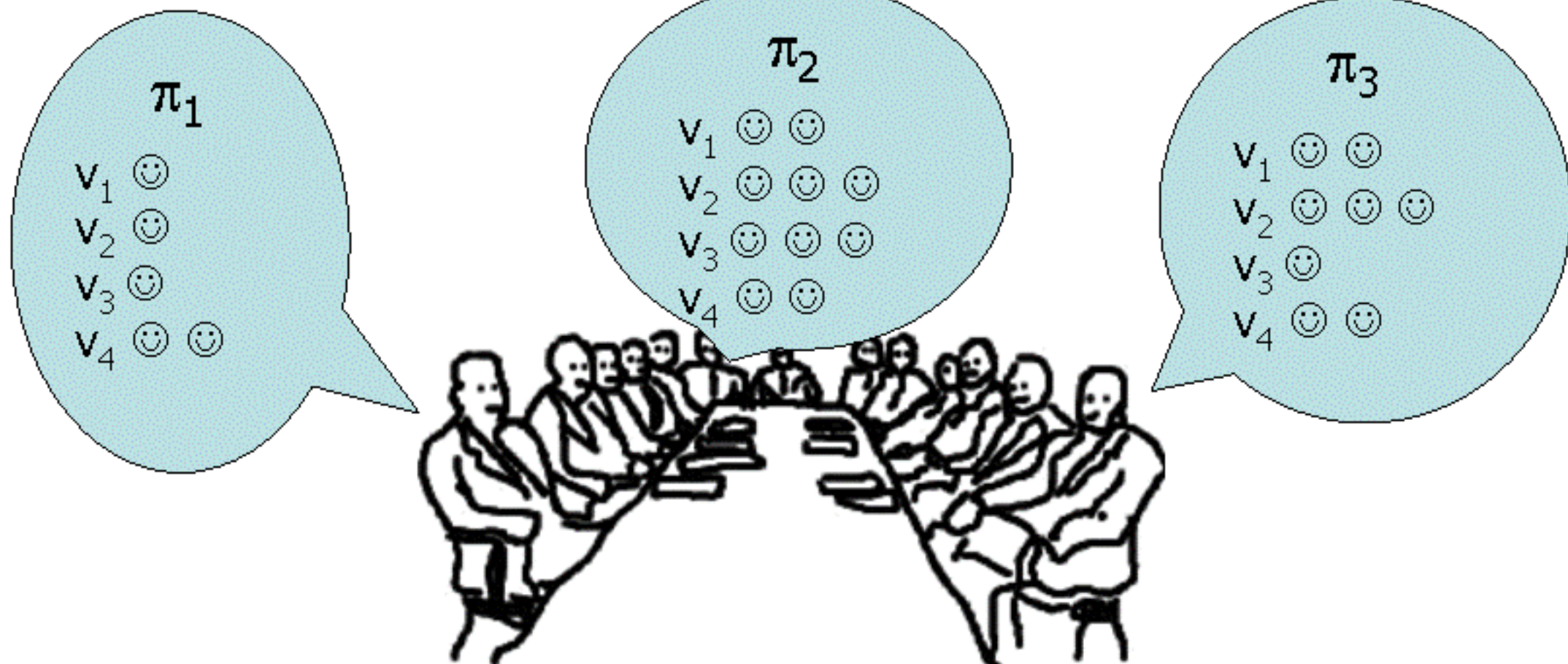
Can pay for ties but “easier” for approximation

## Why minimize over permutations?

Depends on “amount of information” available and  
application

$\langle \cdot, \cdot \rangle$   
not metric on  
 $\Delta \cap \{0,1\}^{v \times v}$   
how to get 2  
approx?





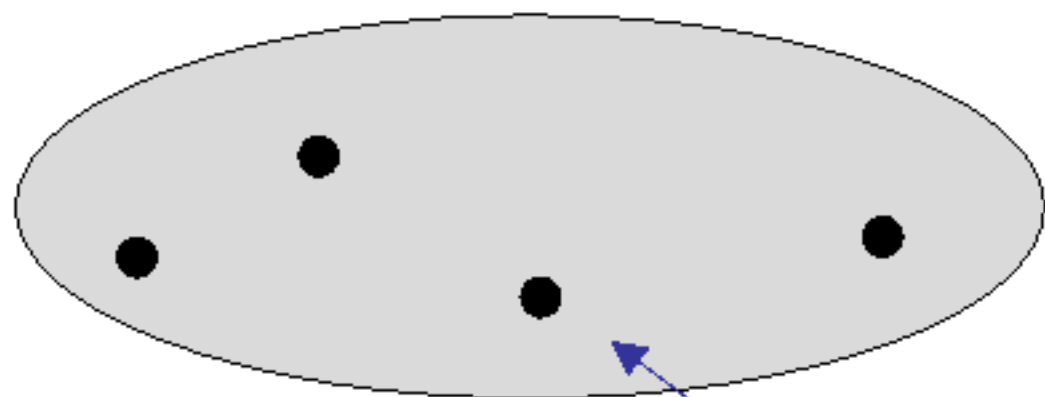
ranking of 4 executives

$$V = \{v_1, v_2, v_3, v_4\}$$

algorithm  
RepeatChoice:  
2 approximation

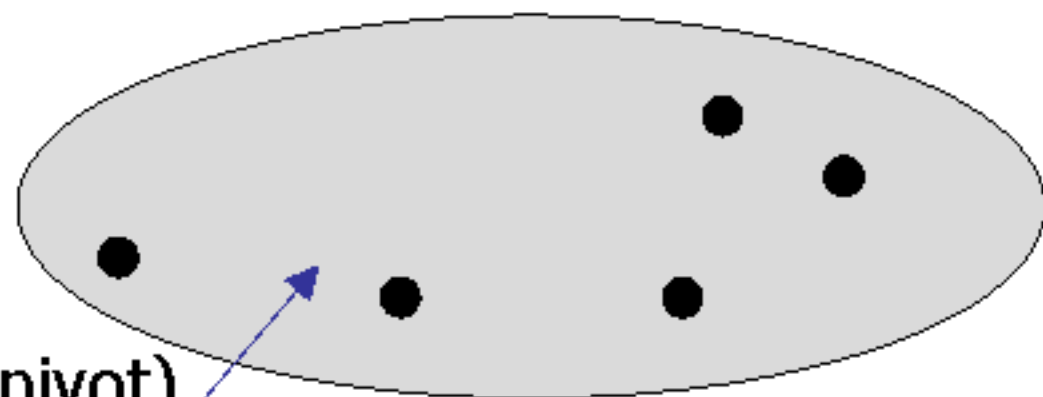
$[v_2, v_3], [v_1, v_4]$  (chose voter #2)  
 $v_2, v_3, [v_1, v_4]$  (broke ties with voter #3)  
 $v_2, v_3, v_4, v_1$  (broke ties with voter #1)

3/2 approximation  
algorithm:



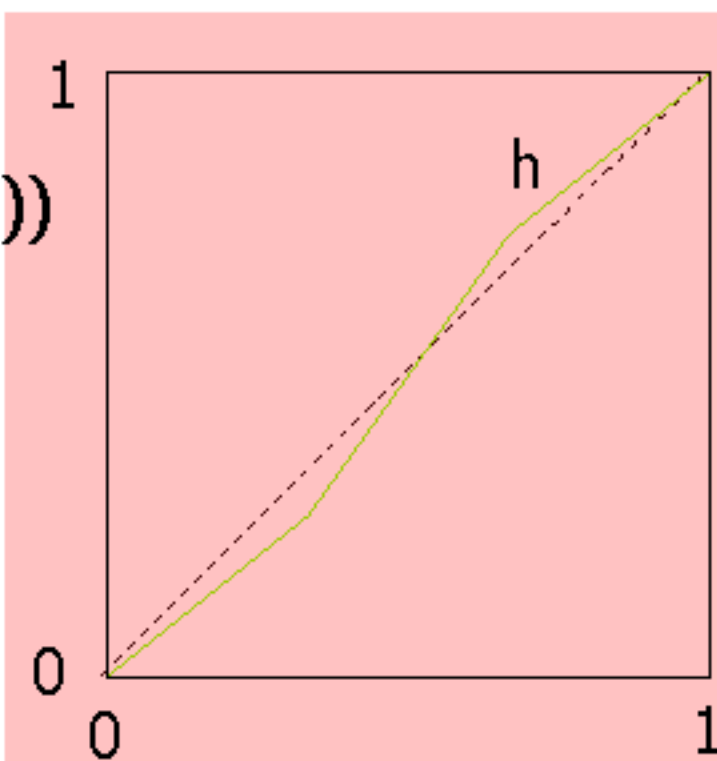
w. pr.  $h(\tau(u,v))$

$u$  (pivot)



w. pr.  $h(\tau(v,u))$

conjecture:  
best of RepeatChoice  
and this gives  
4/3 approximation

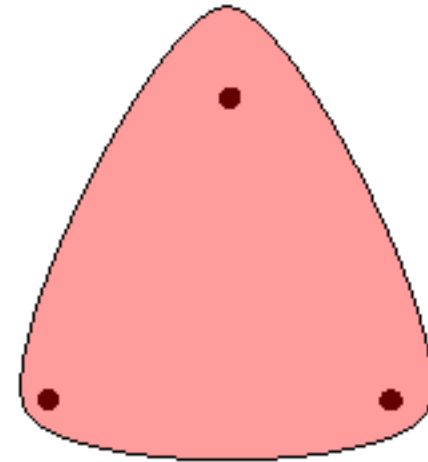
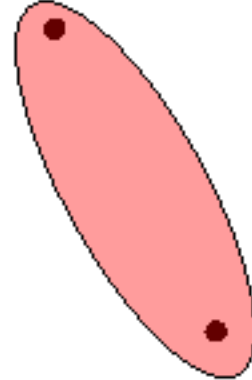
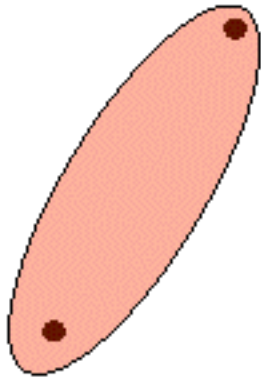


# Followup Work

- Coppersmith, Fleischer, Rudra:  
Borda meets Condorcet
- Hedge, Jain, Williamson, Van Zuylen:  
Derandomization
- Mathieu-Kenyon, Schudy:  
PTAS!

ranking with ties = ranking + clustering

clustering



aggregating . . . .

•  
like ? like  
• dislike •



## Correlation Clustering (on complete graphs):

given set  $V$  of  $n$  vertices,  $w: \binom{V}{2} \rightarrow \mathbb{R}^+$

minimize  $\langle \sigma, w \rangle$  over clustering  $\sigma$

pay  $w\{u,v\}$  whenever  $u \equiv_{\sigma} v$

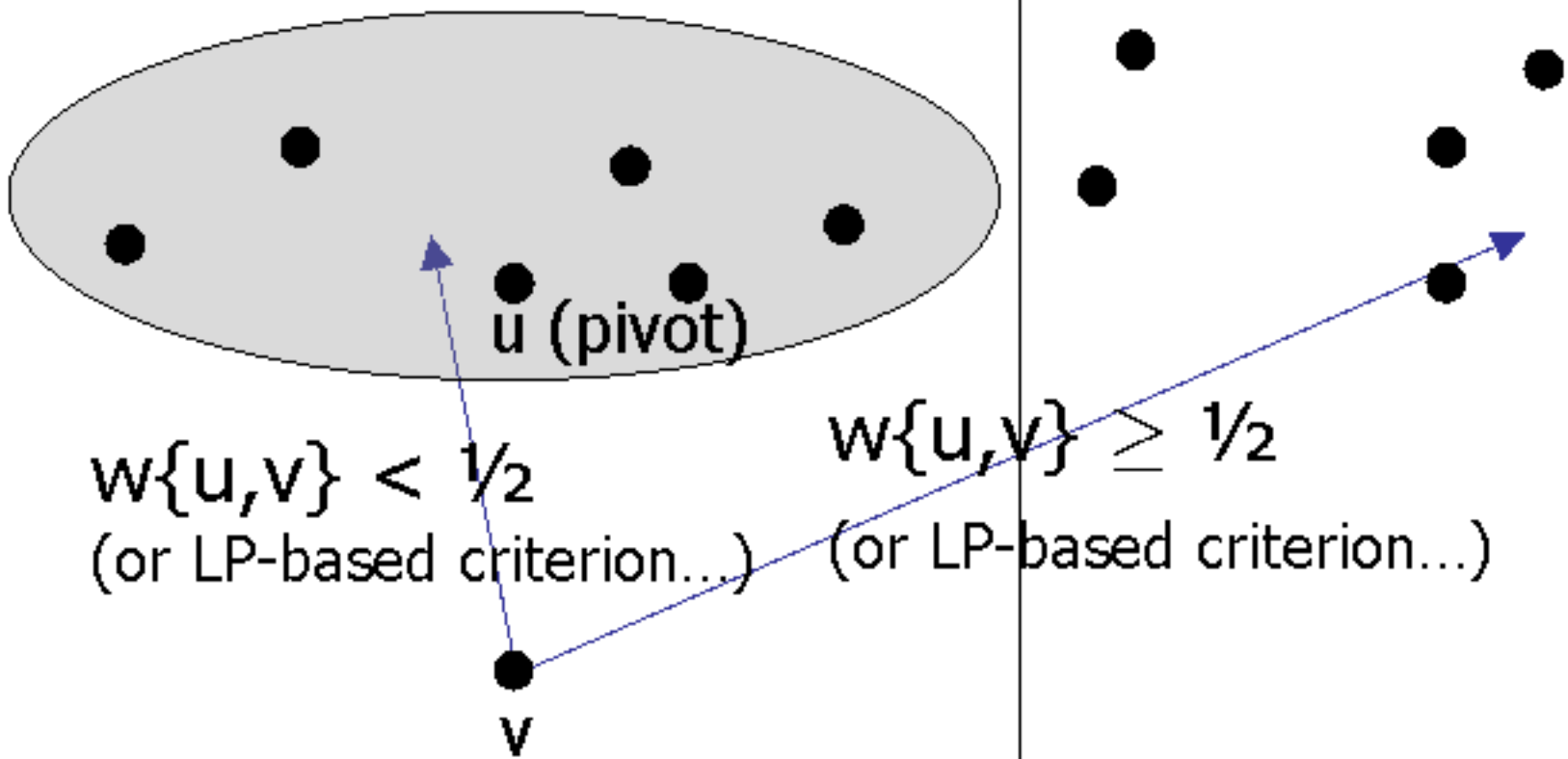
pay  $1-w\{u,v\}$  whenever  $u \not\equiv_{\sigma} v$

## Consensus Clustering:

$w \in \text{conv}(\text{clusterings})$



Time can be quadratic w. prob. 1 for some inputs



### (Combinatorial/LP-based) KwikCluster

## “LP-based” KwikCluster

input	approx	previous	hardness
$[0,1]^{V \times V}$	2.5(corr-clust)	4	Max-SNP-Hard
$\{0,1\}^{V \times V}$	2.5	4	Max-SNP-Hard
$\Delta$	2	4	Max-SNP-Hard
Cons Clust	2	2	Max-SNP-Hard

best of two: 4/3

