

What is a J-L transform?

Def:

A distribution \mathcal{D} on real $k \times d$ matrices has the J-L property if

$$\forall x \in \mathbb{R}^d \quad \Pr_{\phi \sim \mathcal{D}} [\|\phi x\|_2 \geq \epsilon \|x\|_2] \leq e^{-k\epsilon^2} \quad (0 \leq \epsilon \leq 1)$$

Definition not empty (JL '84)

Nontrivial constructions (AC'06)

$\text{TIME}(\mathcal{D}) = \#$ of operations to reduce vector

THM [A. Liberty '07]

For $k > d^{1/2}$ $\exists \mathcal{D}$ with

$$\text{TIME}(\mathcal{D}) = O(d \log d)$$

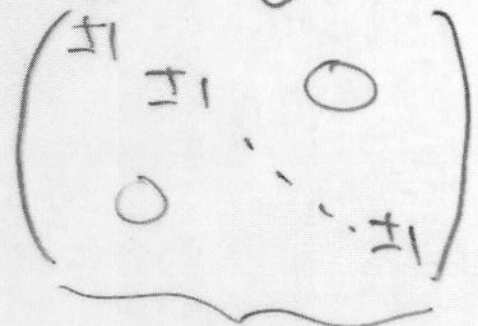
$$k > d^{1/n}$$

THE ALGORITHM

$$r = \frac{1}{8}$$

$\Phi =$

$$B D H D^{(r)} \dots H D^{(r)} H D^{(r)}$$



$d \times d$ Walsh-Hadamard matrix. $O(d \log d)$ -time

$O(d)$ -time

$$H = \frac{1}{\sqrt{d}} \left(\begin{array}{c} \pm 1's \\ \dots \\ \dots \\ \dots \end{array} \right) \Bigg\} d$$

$$B = \frac{1}{\sqrt{k}} \left(\begin{array}{c} \pm 1's \\ \dots \\ \dots \\ \dots \end{array} \right) \Bigg\} k$$

" $B \ll H$ " $\Rightarrow O(d \log d)$ -time
 (rows of B are subset of rows of H)

Consider random variable Z :

$$Z = \left\| \underbrace{B D}_{\substack{\text{matrix} \\ \text{of } \pm 1\text{'s}}} \right\|_2 = \sum_{i=1}^d (B^{(i)} x_i) D_i$$

\uparrow i^{th} column of B \uparrow Random ± 1

$\underbrace{\left(\begin{matrix} \pm 1\text{'s} \\ \vdots \\ \pm 1\text{'s} \end{matrix} \right)}_d \frac{1}{\sqrt{k}}$

Def: $M^{(i)} = B^{(i)} x_i$

$$Z = f(D_1, \dots, D_d) = \left\| \sum M^{(i)} D_i \right\|_2 \in \mathbb{R}$$

$$f: \mathbb{R}^d \rightarrow \mathbb{R}$$

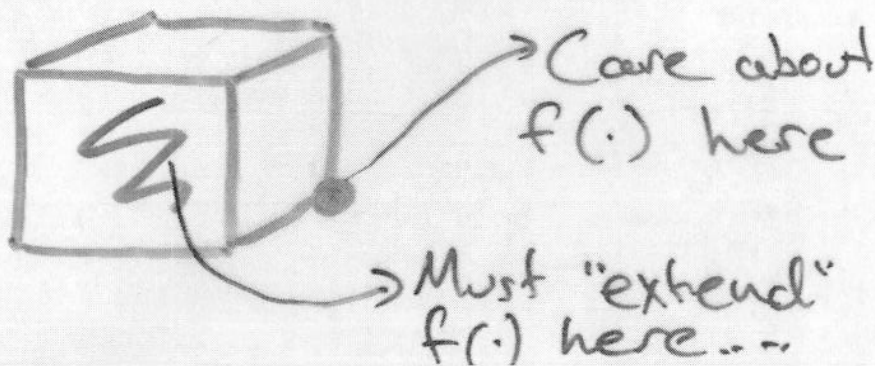
f convex

f σ -Lipschitz

$$\begin{aligned} & \text{for } x, y \in \mathbb{R}^d \\ & |f(x) - f(y)| \leq \sigma \|x - y\|_2 \end{aligned}$$

(Talagrand)

$$\Rightarrow \mathbb{P}_{D_1, \dots, D_d \in \{\pm 1\}} \left[\left| \underbrace{f(D_1, \dots, D_d)}_Z - \underbrace{M_f}_{\text{Median}} \right| > \epsilon \right] \leq e^{-\frac{\epsilon^2}{\sigma^2}}$$



Q: What is σ for f ?

A: It is

$$\sigma = \|M\|_{2 \rightarrow 2}$$

which is

$$\sigma = \|M\|_{2 \rightarrow 2} = \sup_{\|y\|_2=1} \|y^T M\|_2$$

$$= \sup_{\|y\|_2=1} \sqrt{\sum (y^T M^{(i)})^2}$$

$$= \sup_{\|y\|_2=1} \sqrt{\sum x_i^2 (y^T B^{(i)})^2}$$

$$\leq \sup_{\|y\|_2=1} \|x\|_4 \|y^T B\|_4$$

Cauchy
Schwarz

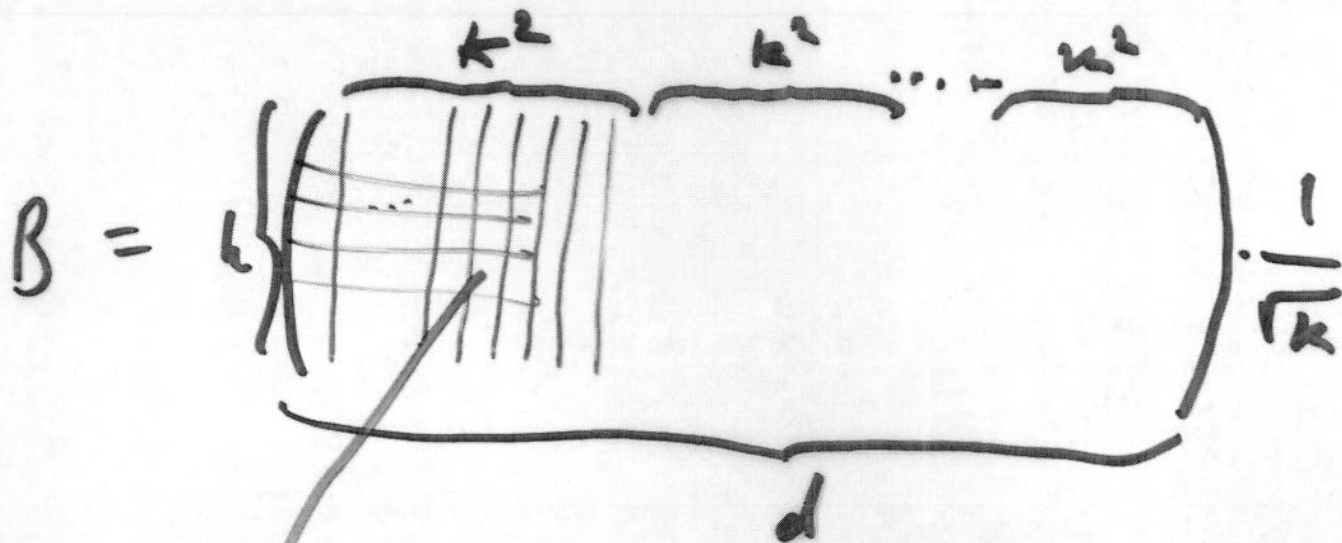
$$= \|x\|_4 \|B^T\|_{2 \rightarrow 4}$$

Q: What is Median of f ?

A: $\approx \|x\|_2 = 1$
(Assumption)

$$P_r[|f - 1| > \epsilon] \leq e^{-\frac{\epsilon^2}{\|x\|_4^2 \|B^T\|_{2 \rightarrow 4}^2}}$$

Constructing B s.t. $\|B\|_2 \rightarrow 4 \frac{d^{1/4}}{k^{1/2}}$



"dual BCH code of designed distance 5 and length k "

- 4-wise independent \Rightarrow Bound
- "BCH" * random $\Rightarrow \|x\|_4 \approx d^{1/4}$

$$P_r[\|f\| > \varepsilon] \leq e^{-\frac{\varepsilon^2 \cdot k}{11 \times 11_4^2 d^{1/2}}}$$

$$\text{get } \|HO^{(r)} \dots HO^{(1)}\|_4 = d^{-1/4}$$

$$P_r \left[\| (BOHO^{(r)}) \dots (HO^{(1)}) \|_2 - 1 > \varepsilon \right] \leq e^{-k\varepsilon^2}$$

($0 \leq \varepsilon \leq 1$)

Controlling $\|x\|_4$

$$Y = \|HOx\|_4$$

$$\sigma \leq \|x\|_4 \|H\|_{\frac{4}{3} \rightarrow 4} = \|x\|_4 d^{-1/4} d^{-\delta/2}$$

$$M_Y = O(d^{-1/4})$$

(median)

$$\|x\|_4 d$$

1

Talagrand:

$$\Pr[\|HOx\|_4 > d^{-1/4} + t] \leq e^{-\frac{t^2}{\|x\|_4^2 d^{-1/2}}}$$

$$t = \|x\|_4 d^{-\delta/2}$$

$$\Pr[\|HOx\|_4 > d^{-1/4} + t] \leq e^{-d^{1/2} - \delta}$$

$$\leq e^{-\kappa}$$

(Assumption)

Repeat... after r times

$$\Pr[\|HO^{(r)} \dots HO^{(1)}x\|_4 > O(d^{-1/4} + \|x\|_4 d^{-\frac{\delta}{2}})]$$

$$\leq r e^{-\kappa}$$

$$\Rightarrow \text{set } r = 1/\delta \dots$$

Lower Bounds

Any ^{non-trivial} lower bound for
say, $k = d^{1/2}$ would
imply lower bound for
FFT which is a
major open problem.