

Distributed Localization and Clustering Using Data Correlation and the Occam's Razor Principle

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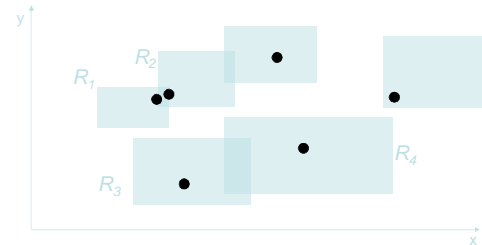
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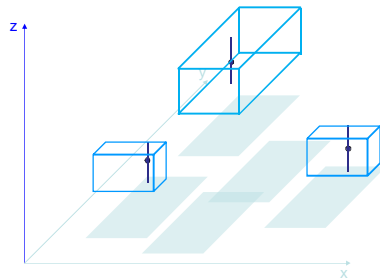
n sensors

- Location of s is not known exactly
 - region R_s in which it lies



Measurements

- z_s
- error e_s

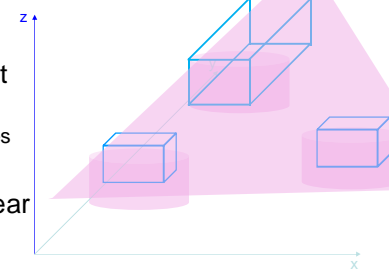


Regression

Measurement = $f(\text{location})$

$z = f(x,y)$

f - consistent
(graph of) f
intersect all boxes
-- "simple"
(piecewise-)linear



Occam Razor's Principle

William of Ockham (ca. 1285-1349):

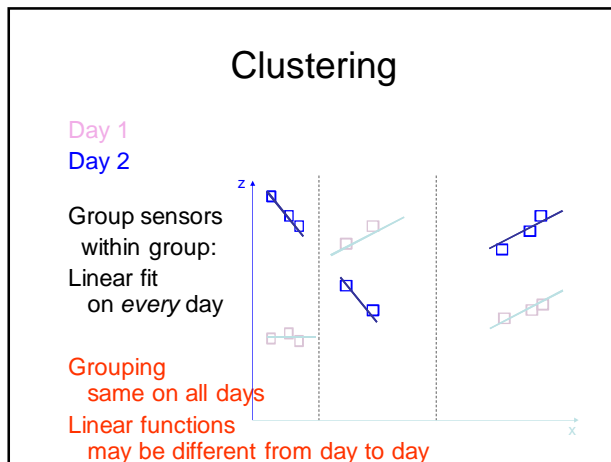
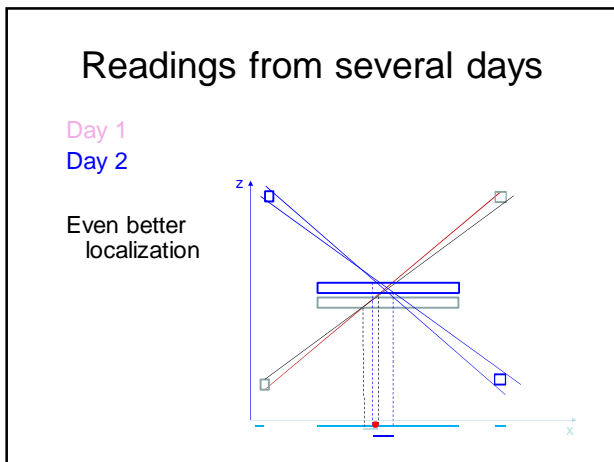
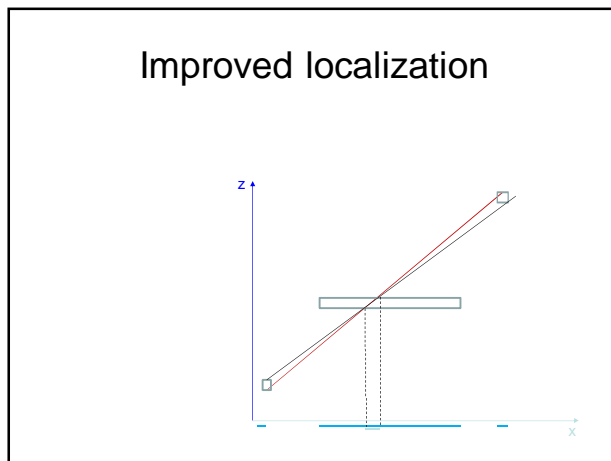
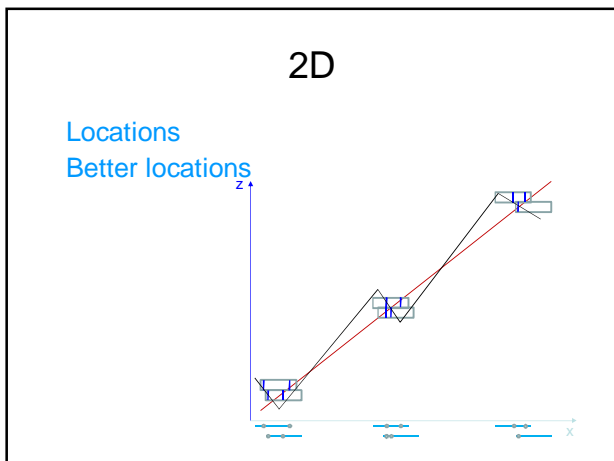
"Entities must not be multiplied beyond necessity"

In other words:

"Given several explanations to a phenomenon, choose the simplest one"

Regression: Know more about z

Our view: Know more also
about (x,y)



Distributed Localization and Clustering Using Data Correlation and the Occam's Razor Principle

3D

- Partition the plane into triangles (Occam: as few as possible) **At the same time: Use linear functions to improve localization**
- Sensors within a triangle: Linear fit on every day

Grouping same on all days
Linear functions may be different from day to day

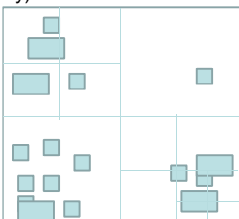
- ### Related work
- Terrain simplification [Agarwal and Desikan'97, Agarwal and Suri'98]
 - NP-hard
 - log OPT - apx
 - Angle-based localization [Basu, Gao, Mitchell, Sabhnani '06]
 - SLAT, dictionary learning, GPS,...

Quad-Tree approach

- Basic problem (oracle):
 - Given a (sub)set S of sensors, find
 - locations p_s in R_s (one per sensor)
 - linear functions h^d (one per day)

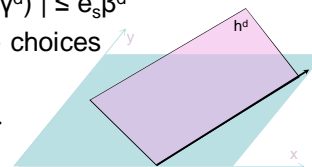
such that $|z_s^d - h^d(p_s)| \leq e_s$

If impossible, split S into 4 quadrants recurse



Basic problem

- Find p_s, h^d s.t. $|z_s^d - h^d(p_s)| \leq e_s$
- $|z_s^d - (a^d x_s + b^d y_s + c^d)| \leq e_s$
 x_s, y_s, a^d, b^d, c^d – variables
- Reparametrizaion $\alpha^d = a^d/b^d, \beta^d = 1/b^d, \gamma^d = c^d/b^d$
 $|z_s^d \beta^d - (\alpha^d x_s + y_s + \gamma^d)| \leq e_s \beta^d$
- α^d : 32 fixed discrete choices
- LP $\rightarrow \{p_s\}$ in $R_s, \{h^d\}$



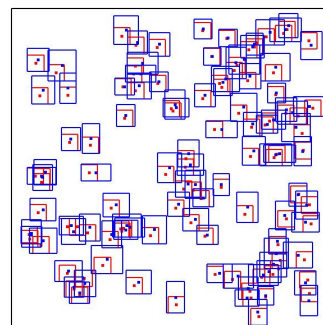
Randomized distributed algorithm

- Find $(x_s, y_s), \beta^d, \gamma^d$ such that $|z_s^d \beta^d - (\alpha^d x_s + y_s + \gamma^d)| \leq e_s \beta^d, (x_s, y_s)$ in R_s
- $Q_s =$ set of $(\beta^1, \gamma^1, \beta^2, \gamma^2, \dots, \beta^D, \gamma^D)$ good for s
- LP is feasible iff $\bigcap Q_s$ is non-empty
- Sensors send to Q_s a leader
 - random set of D^2 sensors do
 - with high probability
 - Terminate in $D \log|S|$ rounds if LP feasible
 - LP is infeasible

K. Clarkson, "Las Vegas algorithms for linear and integer programming when the dimension is small", J. ACM, 1995

Localization

- 1 day
- 100 nodes
- R_s blue
- Red – improved



Clustering

True, original: 4 rooms Output: color-coded clusters

